

The *Energy Spreading* PONS Transform
and its Applications

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<https://prometheus-us.com/PONS-papers/>

Mathematical Features

- Hadamard matrix
- quadrature mirror filter
- low crest factor array; every row of every PONS matrix has the uniform crest factor $\sqrt{2}$
- yields optimal uncertainty principle bounds
- excellent correlation properties

Mathematical Features II

- Each row of every PONS matrix is an energy spreading second order Reed-Muller codeword
- unlike the Walsh-Hadamard matrices, all PONS matrices are symmetric and all even-size PONS matrices have equal row (hence column) sums

Why Spread Energy?

- Covert transmission - signals appear as white noise
- Naturally low probability of intercept, low probability of detection, anti-jam (LPI/LPD/AJ) waveforms
- Robust transmission - gradual signal degradation when corrupted
- Signal can face extreme interference and usable data can still be reconstructed

Computational Behavior

- fast transform: $2\log_2 n$ operations per sample for window size n
- integer computations; no floating point operations
- in-place algorithm
- highly parallelizable
- FPGA built with minimal # of gates
- Algorithms operate to compress and decompress signals in real time

Theorem (H.S. Shapiro — corollary to a compactness result of Kolmogorov):

The following is impossible: $\mathcal{F} \subset L^2(\mathbf{R})$ an infinite orthonormal set, $p > \frac{1}{2}$, and $\forall f \in \mathcal{F}$

$$(*) \quad |f(x)| < \frac{C_1}{(1+|x|)^p}, \quad |\hat{f}(\xi)| < \frac{C_2}{(1+|\xi|)^p}$$

Theorem

\exists a CONS $\{\phi_n\}$ for $C[-\pi, \pi]$ such that

1) $\forall n, x, \phi_n(x)$ takes on only ± 1

$$2) \quad \left| \int_{-\pi}^{\pi} \phi_n(x) e^{-ix\xi} dx \right| \leq \frac{18}{\sqrt{\pi} \sqrt{1+|\xi|}}$$

Corollary (a Global Uncertainty Principle)

\exists a CONS \mathcal{S} for $L^2(\mathbf{R})$ such that (*) with

$p = \frac{1}{2}$ is satisfied for all $f \in \mathcal{S}$

Shapiro Polynomials

$$P_0(z) = Q_0(z) = 1$$

$$P_{n+1}(z) = P_n(z) + z^{2^n} Q_n(z)$$

$$Q_{n+1}(z) = P_n(z) - z^{2^n} Q_n(z)$$

P_n and Q_n are polynomials of degree $2^n - 1$ with coefficients ± 1 .

P_0				1
Q_0				1
P_1		1	1	
Q_1		1	-1	
P_2	1	1	1	-1
Q_2	1	1	-1	1
			⋮	

Properties of Shapiro Polynomials

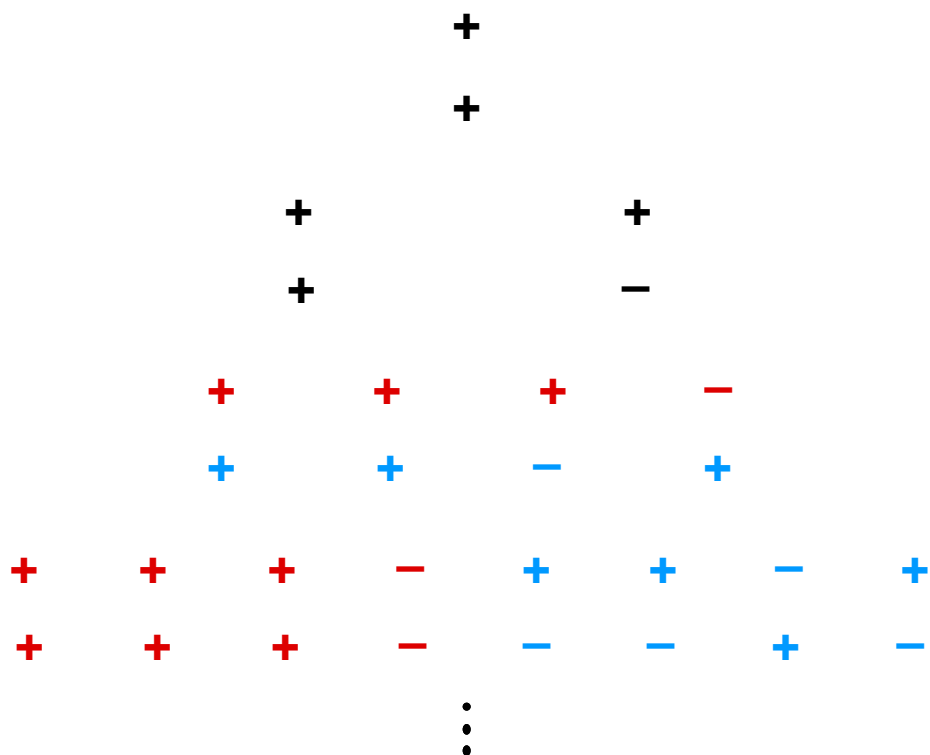
$$\begin{aligned}\text{For } |z|=1, \quad & |P_{n+1}(z)|^2 + |Q_{n+1}(z)|^2 \\ &= 2(|P_n(z)|^2 + |Q_n(z)|^2) \\ &= 2^{n+2}\end{aligned}$$

$$|P_n(z)| \leq \sqrt{2} \sqrt{2^n} \quad \text{or} \quad \|P_n\|_\infty \leq \sqrt{2} \|P_n\|_{L^2}$$

(same for Q_n)

This choice of ± 1 's gives an excellent bound ($\sqrt{2}$) for the “peak factor” (peak-to-average ratio), thereby spreading the “energy” of these polynomials almost equally around the unit circle.

Shapiro Sequences



Three ways of thinking of these sequences:

- (a) sequences of ± 1 's of length 2^n
- (b) coefficients of polynomials
- (c) values of piecewise constant functions on $[-\pi, \pi]$

Note: P_n and Q_n are orthogonal in the sense of (c)

Shapiro Sequences are Incomplete

Want:

A collection of functions of type (c) which form a basis for length 2^n digital signals.

Have:

2 such functions for each n

Need:

2^n such functions for each n

PONS Sequences

+	+
+	-

+	+	+	-
+	+	-	+
+	-	+	+
-	+	+	+

+	+	+	-	+	+	-	+
+	+	+	-	-	-	+	-
+	+	-	+	+	+	+	-
-	-	+	-	+	+	+	-
+	-	+	+	-	+	+	+
+	-	+	+	+	-	-	-
-	+	+	+	+	-	+	+
+	-	-	-	+	-	+	+

⋮

Mathematical Properties

$$P(z) = \sum_{k=0}^{2^n-1} \varepsilon_k z^k, \quad z = e^{it}, \quad \{\varepsilon_k\} \text{ a PONS sequence}$$

$$|P(z)| \leq \sqrt{2} \sqrt{2^n}$$

$$P(z)P(1/z) + P(-z)P(-1/z) \equiv 2^{n+1}$$

$$\left| \int_{-\pi}^{\pi} P(e^{it}) e^{-i\omega t} dt \right| \leq \frac{18}{\sqrt{\pi} \sqrt{1+|\omega|}}, \quad \omega \in \mathbf{R}$$

Walsh Functions

The Walsh functions are

- piecewise constant
- take on only the values ± 1
- form a basis

These properties make the Walsh functions useful in

- signal processing
- digital filtering
- communications
- coding

This usefulness is hampered because

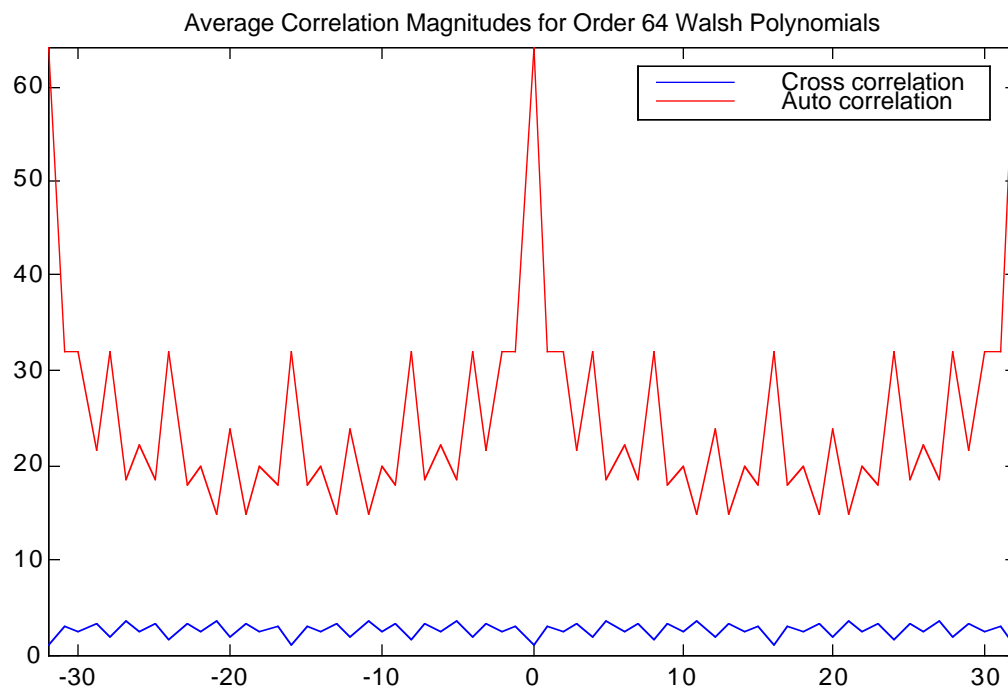
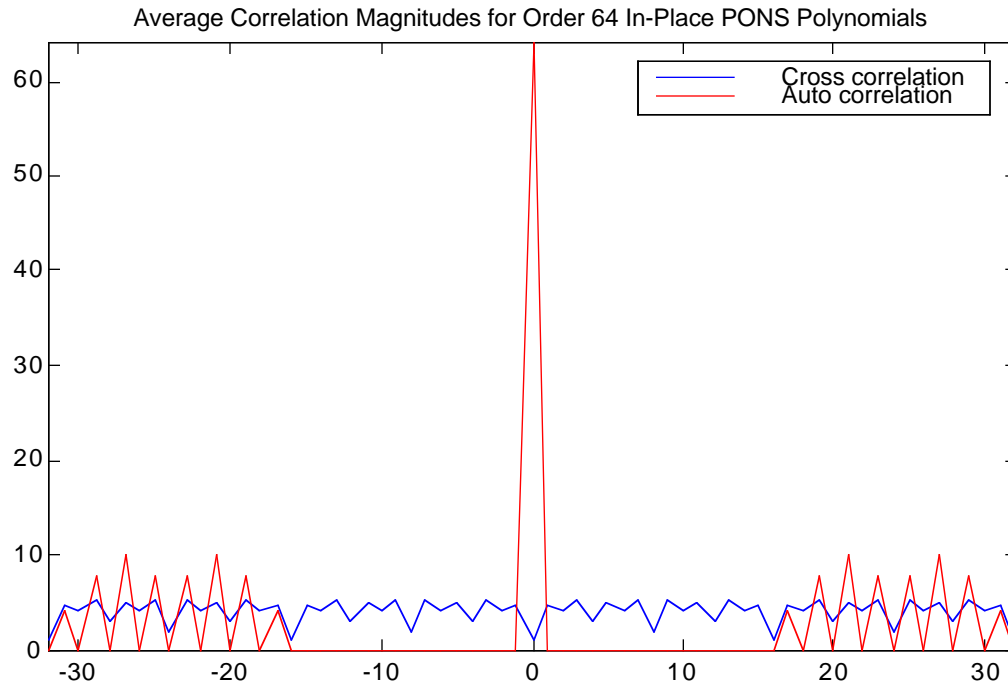
- Walsh polynomials have the worst possible crest factor
- they are not quadrature mirror filters
- their Fourier transforms decay as slowly as possible

PONS = Walsh+

PONS satisfies all useful Walsh properties *plus*

- PONS polynomials have uniformly low crest factor
- they *are* quadrature mirror filters
- PONS Fourier transforms decay as *quickly* as possible
- they are optimal w.r.t. the *Global Uncertainty Principle*

Average Correlation Magnitudes



Energy Spreading in Mathematical Terms

A an $N \times N$ Hadamard matrix, x a discrete signal of length N , $\|\bullet\|_\infty$ the sup norm, $\|\bullet\|_{ME}$ the Beurling minimal extrapolation norm, S a set of such x 's.

Trivial bound : $\|Ax\|_\infty \leq \sqrt{N}\|x\|_\infty$ for all x

Claim : If, for some constant M not too much larger than 1, $\|Ax\|_\infty \leq M\|x\|_\infty$ for all $x \in S$, then A gives good energy spreading on S .

Energy Spreading — Theorems

Theorem: If A is a PONS matrix, then

$$\|Ax\|_{\infty} \leq \sqrt{2}\|x\|_{ME} \text{ for all } x.$$

Example: $x = [x' \mid x'']$, $k + l = N$, A PONS

$$x' = [a \cos(\omega_1 + b), \dots, a \cos(k\omega_1 + b)]$$

$$x'' = [c \cos(\omega_2 + d), \dots, c \cos(l\omega_2 + d)]$$

$$\text{Then } \|Ax\|_{\infty} \leq 3\sqrt{6/5} \sqrt{a^2 + c^2}$$